

# SET THEORY

2020-12-14

① THE WEAK COMPACTNESS THM

② INDSCRIBABILITY.

THE INFINITARY LANGUAGES ALLOW YOU TO CHEAT.

"THE GROUP G IS CYCLIC"

$$\exists x \forall y \exists m \in \mathbb{Z} : y = x^m$$

NOTING

$$\exists x \forall y \left( \begin{array}{l} y = x \vee y = x * x \vee y = x * x^{-1} \\ \vee y = x * x * x \vee y = x^{-1} * x^{-1} * x^{-1} \\ \vee y = x^{-1} * x^{-1} \vee \dots \end{array} \right)$$

$\mathcal{L}_{\omega, \omega}$

$$\exists_{n < \omega} x_n$$

$$\forall_{n < \omega} x_n$$

$\mathbb{C}_p$

$\mathbb{C}_\pi$

IF  $\mathcal{L}_{\omega, \omega}$  SATISFIES WEAK COMPACTNESS AND  $\kappa$  UNACCESSIBLE THEN  $\kappa$  IS WEAKLY COMPACT

WE PROVED THE TREE-PROPERTY

IF  $\kappa$  IS WEAKLY COMPACT THEN  $L_{\kappa, \kappa}$  SATISFIES WEAK COMPACTNESS.

WE HAVE A LANGUAGE  $\mathcal{L}$  OF TYPE  $L_{\kappa, \kappa}$   
 $\exists_{\beta < \alpha} \forall_{\gamma < \alpha}$   
 $\forall_{\beta < \alpha} \bigwedge_{\gamma < \alpha} \varphi$

SENTENCES

- $\Sigma$  A SET OF FORMULAS  $|\Sigma| = \kappa$
  - EVERY  $S \in \Sigma^{\kappa}$  HAS A MODEL
  - WE MUST CONSTRUCT A MODEL FOR  $\Sigma$
  - CONSIDER ONLY THE SYMBOLS OF  $\mathcal{L}$  THAT OCCUR IN  $\Sigma$  AND DISREGARD THE REST NOW  $|\mathcal{L}| = \kappa$
- $L_0 = \mathcal{L}$

$\beta < \kappa$ : FOR EVERY  $L_\beta$ -FORMULA  $\varphi$  WITH FREE VARIABLES  $\langle v_\gamma : \gamma < \alpha \rangle$  INTRODUCE NEW CONSTANTS  $c(\varphi, \gamma, \beta)$   $\gamma < \alpha$  WE GET  $L_{\beta+1}$

$L_0 \rightarrow L_1 : |L_1| = \kappa$   
 $L_\beta \rightarrow L_{\beta+1} : |L_{\beta+1}| = \kappa$   
 IF  $|L_\beta| = \kappa$

$\beta \leq \gamma < \kappa \implies L_\gamma = \bigcup_{\beta < \gamma} L_\beta$

$L^* = \bigcup_{\beta < \kappa} L_\beta$

IF  $\varphi$  IS AN  $L^*$ -FORMULA  
 WITH FREE VARIABLES  $\langle v_j : j < \alpha \rangle$   
 THEN  $\varphi \in L_\gamma$  FOR SOME (MANY)  $\gamma$   
 BECAUSE  $\varphi$  CONTAINS  
 FEWER THAN  $\kappa$  MANY SYMBOLS  
 $L^*$  HAS CONSTANTS

$\underline{c(\varphi, \gamma, j)}$   $j < \alpha$   
 THAT DO NOT OCCUR IN  $\varphi$  ITSELF.

WE MAKE  $\Sigma$  LARGER  
 FOR EACH  $L^*$ -FORMULA  $\varphi$   
 TAKE  $\gamma_\varphi$  SUCH THAT  $\varphi$  IS  $L_{\gamma_\varphi}$   
 DEFINE

$$\sigma_\varphi := (\exists_{j < \alpha} v_j \varphi(v_j : j < \alpha)) \rightarrow (\varphi(c(\varphi, \gamma_\varphi, j) : j < \alpha))$$

$$\Sigma^* = \Sigma \cup \{ \sigma_\varphi : \varphi \text{ IS AN } L^*\text{-FORMULA} \}$$

$\Sigma^*$  IS LIKE  $\Sigma$  :  
 IF  $S \in [\Sigma^*]^{\leq \kappa}$

THEN  $S$  HAS A MODEL

TAKE MODEL FOR  $S \cap \Sigma$  :  $A$   
 $A$  IS JUST AN  $L$ -STRUCTURE.

MAKE THAT A MODEL FOR  $S$   
 BY INTERPRETING

$$(c(\varphi, \gamma_\varphi, j) : j < \alpha)$$

FOR EVERY  $\varphi$  FOR WHICH

$$\sigma_\varphi \in S \setminus \Sigma$$

IF  $\exists_{j < \alpha} v_j \varphi(v_j : j < \alpha)$  IS TRUE IN  $A$

MAKE SURE  $\varphi(c(\varphi, \gamma_\varphi, j) : j < \alpha)$  IS TRUE

OTHER WISE ARBITRARY

NO CONFLICT

(9)

IF  $\varphi \neq \psi$  THEN

$$\underbrace{\{ \langle \varphi, \beta, \gamma \rangle : \beta \in \alpha \} \cap \{ \langle \varphi, \beta, \gamma \rangle : \beta \in \beta \}} = \emptyset$$

ENUMERATE ALL SENTENCES  
IN  $\mathcal{L}^*$  AS  $\langle \sigma_\alpha : \alpha < \kappa \rangle$

$\mathcal{T}$  WILL BE A SUBTREE OF  $2^{<\kappa}$

$t \in \mathcal{T}$  IF  $\delta = \text{DOM } t \in \kappa$

AND  $t : \delta \rightarrow 2$  IS

SUCH THAT

THERE IS A MODEL  $A$  FOR

$$\Sigma^* \cap \{ \sigma_\alpha : \alpha < \delta \}$$

SUCH THAT

$$\rightarrow A \models \sigma_\alpha \text{ IFF } t(\alpha) = 1$$

$\mathcal{T}$  HAS CARDINALITY  $\kappa$

SO IT HAS A BRANCH  $B$ !

$B$  DETERMINES  $\alpha < \kappa \rightarrow 2$

$$\alpha = \cup B$$

$$\Delta = \{ \sigma_\alpha : \alpha(\alpha) = 1 \}$$

$$\Sigma^* \in \Delta$$

- FOR EVERY  $\sigma \rightarrow \sigma \in \Delta$  OR  
 $\neg \sigma \in \Delta$

AS SOON AS

$$\{ \sigma, \neg \sigma \} \subseteq \{ \sigma_\alpha : \alpha < \delta \}$$

JUST ONE OF  $\sigma$  AND  $\neg \sigma$   
GETS VALUE 1

SO  $\Delta$  IS COMPLETE

WE BUILD A MODEL  $\mathcal{A}$

DOMAIN  $A$  CONSISTS OF ALL CLOSED TERMS OF  $\mathcal{L}^*$

CLOSED TERMS:

CONSTANTS

CLOSED UNDER APPLICATION

OF FUNCTION SYMBOLS

$$\begin{aligned} & \parallel e = ee \quad eee = eeee \\ & \parallel e^{-1} = e^{-1}e = ee^{-1} \end{aligned}$$

### INTERPRETATION

- CONSTANTS  $\longrightarrow$  THEMSELVES
- FUNCTION SYMBOLS  $\longrightarrow$  THE NATURAL THING.

- P PREDICATE:

$$\parallel \mathcal{A} \models P(\tau_1, \dots, \tau_n) \text{ IFF } \underline{P(\tau_1, \dots, \tau_n)} \in \Delta \parallel$$

THE RESULT IS AS IN THE STANDARD PROOF.

FOR ALL SENTENCES  $\sigma$

$$\mathcal{A} \models \sigma \quad \text{IFF} \quad \sigma \in \Delta$$

$$\sigma, \tau \quad \longrightarrow \quad \sigma \vee \tau \quad \sigma \wedge \tau$$

$$\neg \sigma, \quad \sigma \longrightarrow \tau$$

$$\sigma = \underline{\exists \{z \in \mathcal{U}\} \varphi(\mathcal{U}_z : z \in \mathcal{U})} \parallel$$

$\Downarrow$   $\parallel \sigma_\varphi$  IS IN  $\Delta$

IF  $\sigma$  IS IN  $\Delta$  THEN

$$\text{SO IS } \underline{\varphi(\mathcal{C}(\varphi, \{z, \tau_z\} : z \in \mathcal{U}))} \in \Delta$$

$$\parallel \mathcal{A} \models \sigma \parallel$$

IF  $A \models \varphi(\tau_\gamma : \gamma < \alpha)$

THEN  $\varphi(\tau_\gamma : \gamma < \alpha) \in \mathcal{K}$

$\varphi(\tau_\gamma : \gamma < \alpha) \rightarrow \exists_{\gamma < \alpha} \forall_{\gamma} \varphi(\tau_\gamma : \gamma < \alpha)$

SO # MODELS  $\Sigma^*$   
AND HENCE  $\Sigma$

THEOREM [TARSKI-SCOTT]

$\mathcal{K}$  IS WEAKLY COMPACT IFF  
IT IS  $\aleph_1$ -UNDESCRIBABLE

$\Pi_n^0$   $\forall x_0 \exists x_1 \dots \forall x_{n-1} (\varphi \text{ -- (WITHOUT QUANTIFIERS)})$   
 $\Sigma_n^0$   $\exists x_0 \forall x_1 \dots \exists x_{n-1}$

$\Pi_1^1$   $\forall X \varphi(X)$  ← REST FIRST-ORDER WITH  $X$  AS PREDICATE " $x_i \in X$ "  
↑  
SECOND ORDER  
SUBSET OF STRUCTURE

$\Sigma_1^1$   $\exists X \varphi(X)$

$\Pi_1^2$   $\forall \mathcal{X} \varphi(\mathcal{X})$  ← FIRST- AND SECOND ORDER  
↑  
THIRD ORDER  
FAMILY OF SUBSETS

$\kappa$  IS  $\Pi_1^1$ -INDESCRIBABLE MEANS  
WHENEVER  $\sigma$  IS A  $\Pi_1^1$ -SENTENCE

$\forall X \varphi(X)$

← NO FREE VAR'S.

AND  $U \in V_\kappa$

ARE SUCH THAT

$(V_\kappa, \in, U) \models \sigma$

↑ CAN BE USED AS  
A PREDICATE IN  $\sigma$

THEN THERE IS  $\alpha < \kappa$   
SUCH THAT

$(V_\alpha, \in, U \cap V_\alpha) \models \sigma$

$\kappa$  WEAKLY COMPACT

- LET  $C \subseteq \kappa$  BE CUB (ACC AS  $U$ )

-  $\varphi(F)$  SAYS "F IS A FUNCTION"

∧ "IF  $\text{DOM } F$  IS AN ORDINAL  
THEN  $\exists \alpha \forall \beta F(\beta) < \alpha$ "

THIS SAYS " $\kappa$  IS REGULAR"

( $\forall \alpha \exists \beta \sup C \cap \alpha \in C$ )

$(V_\kappa, \in, C) \models \forall F (\varphi(F) \wedge \text{"CUB CUB"})$

THEN  $\exists \alpha < \kappa$

$(V_\alpha, \in, C \cap \alpha) \models \forall F (\varphi(F) \wedge \text{"CUB CUB"})$

↑  
 $\alpha$  IS REGULAR  
 $C \cap \alpha$  IS CUB IN  $\alpha$

HENCE

$\{\alpha < \kappa : \alpha \text{ REGULAR}\}$  IS STATIONARY

" $\kappa$  IS MAHLO"

MAHLO IS  $\Pi_1^1$

$\rightarrow [(\forall X)(X \text{ IS CURS} \rightarrow \exists \alpha \in X (f \alpha = \alpha))$   
"C IS CURS"]

$\} \alpha < \kappa : \alpha \text{ IS MAHLO} \} \text{ IS STATIONARY}$

$\rightarrow$  IF  $\kappa$  IS WEAKLY COMPACT AND  $S \subseteq \kappa$  IS STATIONARY IN  $\kappa$  THEN THERE  $\lambda < \kappa$  SUCH THAT  $S \cap \lambda$  IS STATIONARY IN  $\lambda$  REG. UNCOUNTABLE

-  $\kappa$  IS REGULAR:  $\Pi_1^1$

-  $\kappa$  IS UNCOUNTABLE IS  $\Pi_1^1$  IN  $(V_\kappa, \in)$

-  $S$  IS STATIONARY IS  $\Pi_1^1$  IN  $(V_\kappa, \in, S)$

$(\forall C)(C \text{ IS CURS} \rightarrow \exists \alpha \in C \cap S)$

BY THE MAGIC OF INDESCRIBABILITY

$\exists \lambda < \kappa$

$(V_\lambda, \in, S \cap \lambda) \models \text{"ALL OF THE" ABOVE}$

DEFINITION OF WEAK COMPACTNESS IS

$\Pi_2^1$

$(\forall F)(\exists H) \varphi(F, H)$

YOU KILL THIS IN

$\kappa$  SINGULAR  $\rightarrow \kappa$  IS  $\Pi_m^0$  - DESCRIBABLE FOR SOME  $m$ .

$\exists \underline{f} \exists \underline{\lambda}$   
 $\text{PREN PREN}$  [ "f IS A FUNCTION"  
"DOM f =  $\lambda$ "  
"RAN f IS COFINAL"  
 $\varphi(f, \lambda)$  ]



$$(V_\kappa, \in, \neq, \lambda) \models \varphi(f, \lambda)$$

THERE IS NO  $\alpha < \kappa$  SUCH THAT

$$(V_\alpha, \in, \neq \wedge V_\alpha, \lambda \cap V_\alpha) \models \varphi(f, \lambda)$$

①  $\alpha < \lambda$   $\lambda \cap V_\alpha$  TOO SMALL

②  $\alpha > \lambda$  SOME  $\beta < \lambda$  HAVE NO VALUES.

IF  $\lambda < \kappa \leq 2^\lambda$   $\Pi_1^0$ -DESCRIBIBLE  
 $\lambda$  AND  $f = \mathcal{P}(U) \rightarrow \kappa$

$$\in V_\kappa$$

IF  $\kappa$  IS  $\Pi_1^1$ -INDESCRIBABLE  
 THEN  $\kappa$  IS INACCESSIBLE

WE SHOW  $\kappa$  HAS THE  
 TREE PROPERTY.

TAKE A TREE  $\mathcal{T}$  OF CARDINALITY  $\kappa$   
 WITH SMALL LEVELS

PROOF OF 9.26: WLOG  $\mathcal{T} \subseteq 2^{<\kappa}$

ASSUME  $\mathcal{T}$  HAS NO BRANCH

$$\rightarrow (\forall B)(B: \kappa \rightarrow 2 \Rightarrow (\exists \alpha)(\underline{B \upharpoonright \alpha} \notin \mathcal{T}))$$

HOLDS IN  $(V_\kappa, \in, \mathcal{T})$

IT HOLDS IN  $(V_\lambda, \in, \mathcal{T} \cap 2^{<\lambda})$  FOR  
 SOME  $\lambda < \kappa$

THAT SAYS:  $\mathcal{T}_\lambda = \emptyset$  CONTRADICTION

17.17 IF  $\kappa$  IS WEAKLY COMPACT  
 AND  $U \in V_\kappa$

THEN THERE IS A TRANSITIVE  
 ELEMENTARY EXTENSION

$$(M, \in, U') \text{ OF } (V_\kappa, \in, U)$$

WITH  $\kappa \in M$

$$V = (V_\kappa, \in, U, \{x\}_{x \in V_\kappa})$$

↑  
EACH  $x \in V_\kappa$   
IS A CONSTANT

$\Sigma$ : ALL  $L_{\kappa, \kappa}$ -SENTENCES  
IN OUR LANGUAGE  
 $\in, U, \{x\}_{x \in V_\kappa} \cup \{c\}$   
THAT ARE TRUE IN  $V$   
PLUS ALL THE SENTENCES

→ "IS AN ORIGINAL" " $\alpha \in C$ " ( $\alpha < \kappa$ )

$$|\Sigma| = \kappa$$

$$\text{IF } S \in [\Sigma]^{<\kappa}$$

THEN  $S$  HAS A MODEL  
THIS MODEL IS JUST  $V$   
WITH  $\in, U, x (x \in V_\kappa)$   
INTERPRETED AS THEMSELVES  
TAKE  $C$  LARGE ENOUGH

SO  $\Sigma$  HAS A MODEL!

$$A = (A, E, U_A, \{x_\alpha\}_{\alpha \in V_\kappa}, c_A)$$

REPLACE  $x_\alpha$  BY  $x$

$$\text{SO } V_\kappa \subseteq A$$

$$E \cap V_\kappa^2 = \in$$

$$U_A \cap V_\kappa = U$$

$V \not\subseteq A$  TRUTH OF SENTENCES  
↑  
ELEM  
SUBST

$(A, E)$  IS WELL-FOUNDED

$$\neg (\exists_{n \in \omega} V_n) (\bigwedge_{n \in \omega} V_{n+1} \in V_n)$$

HOLDS IN  $V$   $E$

$L_{\omega_1, \omega_1}$

LET  $\mathcal{M}$  BE THE TRANSITIVE COLLAPSE

THE INTERPRETATION OF  $\mathcal{C}$  IN  $\mathcal{M}$  IS AN ORDINAL ABOVE ALL  $\alpha < \kappa$  SO  $\kappa \in \mathcal{M}$  !!

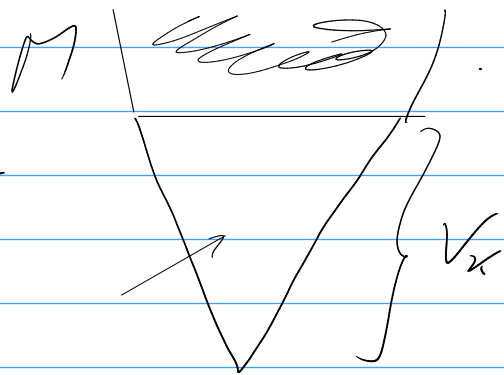
ASSUME  $\kappa$  IS WEAKLY COMPACT  
 $U \in V_\kappa$  AND

(\*)  $(V_{\kappa_1}, \epsilon, U) \models (\forall X) \varphi(X)$

TAKE  $(\mathcal{M}, \epsilon, U')$  AS ABOVE

By ELEMENTARITY AND THE USE OF  $\mathcal{L}_{\kappa, \kappa}$  YOU GET

$V_\kappa^{\mathcal{M}} = V_\kappa$



(\*) MEANS

$(\forall X \in V_\kappa) ((V_{\kappa_1}, \epsilon, U) \models \varphi(X))$

SO THIS ALSO HAPPENS IN  $\mathcal{M}$

$(\mathcal{M}, \epsilon, U') \models (\forall X \in V_\kappa) ((V_{\kappa_1}, \epsilon, U' \cap V_\kappa) \models \varphi(X))$

$\kappa \in \mathcal{M}$

$(\mathcal{M}, \epsilon, U') \models (\exists \alpha) (\forall X \in V_\alpha) ((V_{\alpha_1}, \epsilon, U' \cap V_\alpha) \models \varphi(X))$

ELEMENTARITY

$(V_{\kappa_1}, \epsilon, U) \models \dots$

THIS MEANS  $(\exists \alpha < \kappa) ((V_{\alpha_1}, \epsilon, U \cap V_\alpha) \models \sigma)$ .

ULTRAFILTERS  $\longrightarrow$  MEASURABLE

PARTITIONS /  $\longrightarrow$  WEAKLY  
KÖNIG'S LEMMA COMPACT

REGULAR STRONG LIMIT  $\longrightarrow$  INACCESSIBLE

WEAKLY COMPACT: TONS OF  
INACCESSIBLES BELOW IT

MEASURABLE: TONS OF  
WEAKLY COMPACTS BELOW IT

$$j: V \hookrightarrow M \subset V$$

KUNEN 17.7

THIS IS NECESSARY

THERE IS NO NON-TRIVIAL  
ELEM EMBEDDING  $j: V \rightarrow V$

