

# SET THEORY 2021-12-06

①

TODAY: QUITE TECHNICAL.

- CONSTRUCT  $\mathcal{M}[G]$  FROM  $\mathcal{M}$  AND  $G$ .
- $g: w \rightarrow w_1^\mathcal{M}$      $h: w_1^\mathcal{M} \rightarrow w_2^\mathcal{M}$   
MAKE RELATIONS IN  $\mathcal{M}$   
TO HELP TO SHOW  $g, h$  NOT SURJ.
- $\mathcal{M}[G]$  IS A MODEL OF (ENOUGH OF) ZFC

$\mathcal{P}$  A PARTIAL ORDER IN  $\mathcal{M}$   
 $\text{FN}(w_2^\mathcal{M} \times w, 2)$   
 $\leq$  REFLEXIVE, ANTISYMM, TRANS.  
FOR OUR COMFORT:

$\mathcal{P}$  HAS A MAXIMUM  $\mathbb{1}$   
 $\text{FN}(w_2^\mathcal{M} \times w, 2)$ :  $\mathbb{1} = \emptyset$   
( $p \leq q$  MEANS  $p \supseteq q$ )

$\mathcal{P}$  AND  $q$  ARE COMPATIBLE IF  
THERE IS AN  $r$  WITH  $r \leq p, q$

$\text{FN}(w_2^\mathcal{M} \times w, 2)$      $p \cup q$  IS A FUNCTION  
( $\{x \in \text{DOM } p \cap \text{DOM } q \rightarrow p(x) = q(x)\}$ )  
INCOMPATIBLE IS NOT COMPATIBLE  
NOTATION     $p \perp q$ :  $\exists x \in \text{DOM } p \cap \text{DOM } q$   
 $p(x) \neq q(x)$ .

FILTER ON  $\mathcal{P}$ : SUBSET  $G \neq \emptyset$   
SUCH THAT -  $p, q \in G \rightarrow \exists r \in G$   $r \leq p, q$   
-  $p \in G, q \supseteq p \rightarrow q \in G$   
(SO  $\mathbb{1} \in G$ )

$\mathcal{D} \subseteq \mathcal{P}$  IS DENSE IF  
 $\forall p \in \mathcal{P} \exists q \in \mathcal{D}$   $q \leq p$

A FILTER

$G$  IS  $M$ -GENERIC ON  $P$

IF  $G \cap D \neq \emptyset$  FOR EVERY DENSED THAT IS IN  $M$ .

$\Pi$  COUNTABLE IF  $p \in P$  THEN THERE IS AN  $M$ -GENERIC  $G$  SUCH THAT  $p \in G$ .

$IP$ -NAME: A SET  $\Sigma$  THAT IS A RELATION AND SATISFIES:  
IF  $\langle \pi, p \rangle \in \Sigma$  THEN  $\pi$  IS A  $IP$ -NAME AND  $p \in P$ .

THIS USES RECURSION:

$x \dot{\in} y$  MEANS  $x \in y$  ONLY  $\langle \text{RANK } x < \text{RANK } y \rangle$

$H(\Sigma) = 1$  IFF  $\forall \langle \pi, p \rangle \in \Sigma$   
 $\downarrow$   
 $H(\pi) = 1 \wedge p \in P$   
 $0$  IFF NOT

BEING A NAME IS ABSOLUTE

$\Sigma \in M \quad P \in M$

$\Sigma$  IS A  $IP$ -NAME IFF  $(\Sigma \text{ IS A } IP\text{-NAME})^M$

$\rightarrow \text{VAL}(\Sigma, G) = \{ \text{VAL}(\pi, G) : (\exists p \in G) (\langle \pi, p \rangle \in \Sigma) \}$   
(OR  $\Sigma_G$ )  $\uparrow \emptyset$

$M[G] = \{ \text{VAL}(\Sigma, G) : \Sigma \text{ IS A } IP\text{-NAME IN } M \}$

?  $M \in M[G]$ ?

?  $G \in M[G]$ ?

$M^P$  IS THE CLASS OF  $IP$ -NAMES IN  $M$

FOR  $x \in M$

$\check{x} = \{ \langle \check{y}, 1 \rangle : y \in x \}$  RECURSIVE

$\check{\emptyset} = \emptyset \quad \check{1} = \{ \langle \check{0}, 1 \rangle \} \quad \check{\omega} = \{ \langle \check{n}, 1 \rangle : n \in \omega \}$

$$\begin{aligned} \text{VAL}(\check{x}, G) &= \{ \text{VAL}(\check{y}, G) : y \in x \} \\ &= \{ y : y \in x \} \\ &= x \end{aligned}$$

so  $M \in M[G]$  !

$$\begin{aligned} \Gamma &= \{ \langle \check{p}, p \rangle : p \in P \} \\ \text{VAL}(\Gamma, G) &= \{ \text{VAL}(\check{p}, G) : p \in G \} \\ &= \{ p : p \in G \} \\ &= G \end{aligned}$$

$G \in M[G]$  !

- $M[G]$  IS TRANSITIVE
- $\text{RANK}(\check{\tau}_G) \in \text{RANK}(\tau)$  BY INDUCTION  
SO  $ON \cap M[G] = ON \cap M$
- $\sigma, \tau$  NAMES :  $\{ \langle \sigma, \pi \rangle, \langle \tau, \pi \rangle \} = \text{up}(\sigma, \tau)$   
IS A NAME.

$$\text{VAL}(\text{up}(\sigma, \tau), G) = \{ \sigma_G, \tau_G \}$$

$M[G]$  SATISFIES PAIRING.

$$\begin{aligned} \text{OP}(\sigma, \tau) &= \text{up}(\text{up}(\sigma, \sigma), \text{up}(\sigma, \tau)) \\ &\hookrightarrow \langle \sigma_G, \tau_G \rangle \end{aligned}$$

TRANSITIVITY : EXTENSIONALITY  
REGULARITY

UNION; GIVEN  $\tau$  LET

$$\begin{aligned} \sigma &= \bigcup \text{DOM } \tau \\ &= \bigcup \{ \pi : (\exists p) \langle \pi, p \rangle \in \tau \} \end{aligned}$$

IN  $M[G]$  :

$$\bigcup \tau_G \in \sigma_G \quad \text{CHECK THIS!}$$

GI : BUILD A  $\sigma$  SUCH THAT  
EVEN  $\bigcup \tau_G = \sigma_G$ .

WE HAVE  $\aleph$  INFINITY TOO!

$$\omega = \text{VAL}(\check{\omega}, G) \in M[G].$$

$$\begin{aligned} \text{RANK } x &= \sup \{ \text{RANK } y + 1 : y \in x \} \\ \aleph \in ON \cap M &\rightarrow \underline{V_\aleph^M = V_\aleph \cap M} \end{aligned}$$

# HENRI CARTAN 1937

1. Soit  $\mathcal{E}$  un ensemble donné une fois pour toutes. Une famille  $\mathbf{F}$  de sous-ensembles de  $\mathcal{E}$  prend le nom de *filtre* (construit sur  $\mathcal{E}$ ) si elle remplit les trois conditions suivantes :

F-I : la famille  $\mathbf{F}$  n'est pas vide et ne contient pas le sous-ensemble vide;

F-II : l'intersection de deux ensembles de  $\mathbf{F}$  appartient à  $\mathbf{F}$ ;

F-III : tout ensemble qui contient un ensemble de  $\mathbf{F}$  appartient à  $\mathbf{F}$ .

## FORCING !

WORK WITH  $\mathcal{F}M(\omega_2 \times \omega, 2)$

IS UG IN  $\mathcal{M}[\mathcal{E}]$  ? ?

$$\Psi = \left\{ \left\langle \underbrace{\langle \langle \alpha, n \rangle, i \rangle^V}_{\text{THE 'CHECK' OF THAT ORDERED PAIR}}, p \right\rangle : \begin{array}{l} \langle \alpha, n \rangle \in \text{DOM } p \\ p(\alpha, n) = i \end{array} \right\}$$

THE 'CHECK' OF THAT ORDERED PAIR

$$\Psi_G = \text{UG}$$

MAKE A NAME FOR  $\overline{\Phi}$

$$f : \omega_2 \rightarrow \mathcal{P}(\omega)$$

$$\alpha \mapsto \{ n : \text{UG}(\alpha, n) = 1 \}$$

MEMBERS OF  $\mathcal{P}$  : CONDITIONS

$$q = \{ \langle \langle \alpha, 10 \rangle, 1 \rangle, \langle \langle \alpha+1, 10 \rangle, 0 \rangle \}$$

IF  $q \in G$  THEN  $10 \in \underbrace{f(\alpha)}_{\equiv} \setminus \underbrace{f(\alpha+1)}_{\equiv}$  IN  $\mathcal{M}[\mathcal{E}]$

$$q \Vdash 10 \in \underbrace{\overline{\Phi}(\check{\alpha})}_{\uparrow} \setminus \underbrace{\overline{\Phi}(\check{\alpha}+1)}_{\uparrow} \quad f = \overline{\Phi}_G$$

DEFINITION OF  $\Vdash$

IF  $\varphi$  IS A FORMULA AND

$\tau_1, \dots, \tau_k$  ARE NAMES

THEN  $p \Vdash \varphi(\tau_1, \dots, \tau_k)$  "p FORCES --"

MEANS FOR ALL  $\mathcal{P}$ -GENERIC  $G$  ON  $\mathcal{P}$

WITH  $p \in G$  WE HAVE

$$\varphi(\tau_{1,G}, \dots, \tau_{k,G})^{\mathcal{M}[\mathcal{E}]}$$

$$\mathcal{M}[\mathcal{E}] \models \varphi(\tau_{1,G}, \dots, \tau_{k,G})$$

REMEMBER  $g: \omega \rightarrow \omega_1^M$   
IT HAS A NAME  $\check{\gamma}$

$$g = \check{\gamma}_G$$

$$\underline{M} \ni R_g = \{ \langle p, \langle n, \alpha \rangle \rangle : n \in \omega, \alpha \in \omega, p \in P, \underline{p \Vdash \check{\gamma}(n) = \alpha} \}$$

- ① IS  $R_g$  IN  $M$ ?
- ② DOES IT WORK?  $R_g[G]$  A FUNCTION?
- ① THAT IS OUR NEXT JOB
- ② YES --- ALMOST

WE DEFINE  $\Vdash^*$   
WITHOUT MENTIONING  $G$ 'S  
PROVE  $p \Vdash \varphi$  IFF  $(p \Vdash^* \varphi)^M$

IF THIS WORKS THEN

$$R_g = \{ \langle p, \langle n, \alpha \rangle \rangle : n \in \omega, \alpha \in \omega, p \in P, \underline{p \Vdash^* \check{\gamma}(n) = \alpha} \}$$

IS DEFINED IN  $M$

$$p \Vdash^* \tau_1 = \tau_2$$

(A) FOR ALL  $\langle \pi_1, s_1 \rangle \in \tau_1$

$$\{ q \leq p : q \leq s_1 \rightarrow$$

$$(\exists \langle \pi_2, s_2 \rangle \in \tau_2) (q \leq s_2 \wedge$$

$$\underline{\text{ABSOLUTE}} \quad \underline{q \Vdash^* \pi_1 = \pi_2}) \}$$

IS DENSE BELOW  $p$ .

$$[ \forall r \leq p \exists q \leq r q \in \dots ]$$

(B) INTERCHANGE 1 AND 2

RECURSION ON  $\langle \pi_1, \pi_2 \rangle \in \langle \tau_1, \tau_2 \rangle$   
:  $\pi_1 \in \text{DOM } \tau_1, \pi_2 \in \text{DOM } \tau_2$

$$p \Vdash^* \tau_1 \in \tau_2$$

MEANS  $\{q \leq p\}$

$$\left\{ \exists \langle \pi_1, \pi_2 \rangle \in \tau_2 (q \leq \pi_1 \wedge q \Vdash^* \pi_2 = \tau_1) \right\}$$

IS DENSE BELOW  $p$

CONJUNCTION!

$$\underline{p \Vdash^* \varphi \wedge \psi} \quad \text{IFF} \quad \underline{p \Vdash^* \varphi} \quad \text{AND} \quad \underline{p \Vdash^* \psi}$$

NEGATION:

$$\underline{p \Vdash^* \neg \varphi} \quad \{q \leq p : q \Vdash^* \varphi\} = \emptyset$$

QUANTIFICATION

$$p \Vdash^* (\exists x) \varphi(x, \tau_1, \dots, \tau_n)$$

IFF

$$\{q \leq p : (\exists \sigma) (q \Vdash^* \varphi(\sigma, \tau_1, \dots, \tau_n))\}$$

IS DENSE BELOW  $p$ .

WE WANT

$$\varphi(\tau_{1,G}, \dots, \tau_{n,G}) \stackrel{\text{MEG}}{}$$

IFF

$$(\exists p \in G) (p \Vdash^* \varphi(\tau_1, \dots, \tau_n))$$

## BIG THEOREM

$\mathbb{M}$  TRANSITIVE MODEL OF ZFC

$P$  A PARTIAL ORDER IN  $\mathbb{M}$

$\tau_1, \dots, \tau_n \in \mathbb{M}^P$

$G$  AN  $\mathbb{M}$ -GENERIC FILTER ON  $P$

$\varphi(x_1, \dots, x_n)$  FORMULA

① IF  $p \in G$  AND  $(p \Vdash^* \varphi(\tau_1, \dots, \tau_n)) \stackrel{\mathbb{M}}{}$   
 THEN  $\varphi(\tau_{1,G}, \dots, \tau_{n,G}) \stackrel{\text{MEG}}{}$

② IF  $\varphi(\tau_{1,G}, \dots, \tau_{n,G}) \stackrel{\text{MEG}}{}$   
 THEN THERE IS A  $p \in G$  SUCH THAT  
 $p \Vdash^* \varphi(\tau_1, \dots, \tau_n)$

Full proof in Kunen's Book 1980

JUST FOR  $\tau_1 = \tau_2$

①  $p \Vdash^* \tau_1 = \tau_2$  AND  $p \in G$

TO SHOW  $\tau_{1G} \in \tau_{2G}$  AND  $\tau_{2G} \in \tau_{1G}$ .

TAKE  $\langle \pi_1, s_1 \rangle \in \tau_1$  WITH  $s_1 \in G$

NEED TO SHOW  $\overline{\pi_{1G}} \in \tau_{2G}$

TAKE  $r \in G$  WITH  $r \leq p, s_1$   $r \Vdash^* \tau_1 = \tau_2$

$D = \{ q : q \perp r \text{ OR } q \in s_1 \rightarrow \exists \langle \pi_2, s_2 \rangle \in \tau_2$   
 $q \in s_2 \wedge q \Vdash^* \pi_1 = \pi_2 \}$

TAKE  $q \in G \cap D$ : WLOG  $q \in s_1$

SO THERE IS  $\langle \pi_2, s_2 \rangle \in \tau_2$

WITH  $q \in s_2$  AND  $q \Vdash^* \pi_1 = \pi_2$

INDUCTION  $\overline{\pi_{1G}} = \pi_{2G}$

SO  $\overline{\pi_{1G}} \in \tau_{2G}$

②  $D$  IS THE UNION OF

$D_1 = \{ r \in P : r \Vdash^* \tau_1 = \tau_2 \}$

$D_2 = \{ r \in P : (\exists \langle \pi_1, s_1 \rangle \in \tau_1) (r \leq s_1 \wedge$   
 $(\forall \langle \pi_2, s_2 \rangle \in \tau_2) (\forall q \in P)$   
 $(q \in s_2 \wedge q \Vdash^* \pi_1 = \pi_2 \rightarrow q \perp r)) \}$

$r \in G : \overline{\pi_{1G}} \in \tau_{1G}$

IF  $q$  THINKS  $\overline{\pi_{2G}} \in \tau_{2G}$

$q$  WOULD GIVE  $\overline{\pi_{2G}} = \overline{\pi_{1G}}$

THEN  $q \in G$

$D_3$  LIKE  $D_2$  BUT WITH 1 AND 2 INTERCHANGED

$D_1 \cup D_2 \cup D_3$  IS DENSE

IF  $p \in P$  THEN  $p \Vdash^* \tau_1 = \tau_2 : p \in D_1$

OR NOT: THEN THERE

IS A  $\langle \pi_1, s_1 \rangle$  WHERE

$\{ q : q \leq s_1 \rightarrow \exists \langle \pi_2, s_2 \rangle \in \tau_2$   
 $q \in s_2 \wedge q \Vdash^* \pi_1 = \pi_2 \}$

IS NOT DENSE BELOW  $p$

WE HAVE  $\mathcal{Z} \in \mathcal{P}$  WITH NO SUCH  $q$  BELOW  $\mathcal{Z}$

$$\left. \begin{aligned} \forall q \in \mathcal{Z} \quad q \in S_1 \wedge \forall \langle \pi_1, s_2 \rangle \in \mathcal{C} \\ (q \notin S_2 \vee q \nVdash^* \pi_1 = \pi_2) \end{aligned} \right\}$$

NOW USE CONTRA POSITIVE

SO WE FIND  $G \cap (D_1 \cup D_2 \cup D_3) \neq \emptyset$

IF  $p \in G \cap D_2$  THEN WE

FIND  $\langle \pi_1, s_1 \rangle \in \mathcal{C}_1$  WITH  $p \in S_1$

SO  $s_1 \in G$  AND  $\pi_1 \in \mathcal{C}_1$

IF  $\pi_1 = \pi_2$  FOR SOME  $\langle \pi_2, s_2 \rangle \in \mathcal{C}_2$

WITH  $s_2 \in G$

THERE IS  $q_0 \in G$   $q_0 \nVdash^* \pi_1 = \pi_2$

$q_0 \in p_1, s_1, s_2$  SO  $q_0 \perp p$  CONTRADICTION

SO  $G \cap D_1 \neq \emptyset$  AND WE ARE DONE

COROLLARY

$$\rightarrow \underline{p \nVdash \varphi(\mathcal{C}_1, \dots, \mathcal{C}_n)} \text{ IFF } \underline{(p \nVdash^* \varphi(\mathcal{C}_1, \dots, \mathcal{C}_n))}^{\text{M}}$$

FOR ALL  $G$  WE GET

$$\varphi^{[G]} \text{ IFF } (\exists p \in G)(p \nVdash \varphi)$$

WE TAKE  $g: W \rightarrow W_1^M$

WE HAVE  $f \in M^P: g = f \circ \pi$

( $g$  IS A FUNCTION FROM  $W$  TO  $W_1^M$ )<sup>[G]</sup>

THERE IS A  $p \in G$  SUCH THAT

$$(p \nVdash^* f \text{ IS A FUNCTION FROM } W \text{ TO } W_1^M)^{\text{M}}$$

Rg TWO PARTS

$$\textcircled{1} \quad \{ \langle q, \langle \pi, \alpha \rangle \rangle : q \in p, \underline{q \nVdash f(\check{\pi}) = \check{\alpha}} \}$$

$$\textcircled{2} \quad \{ \langle q, \langle \pi, 0 \rangle \rangle : q \perp p, \text{ NEW} \}$$

IF  $p \in H$  THEN  $\delta_H$  IS A FUNCTION

IF  $p \in H$  THEN  $\delta_H$  IS THE ZERO FUNCTION.











